



IBM Thomas J. Watson Research Center

# Parallel-Server Stochastic Systems with Dynamic Affinity Scheduling and Load Balancing

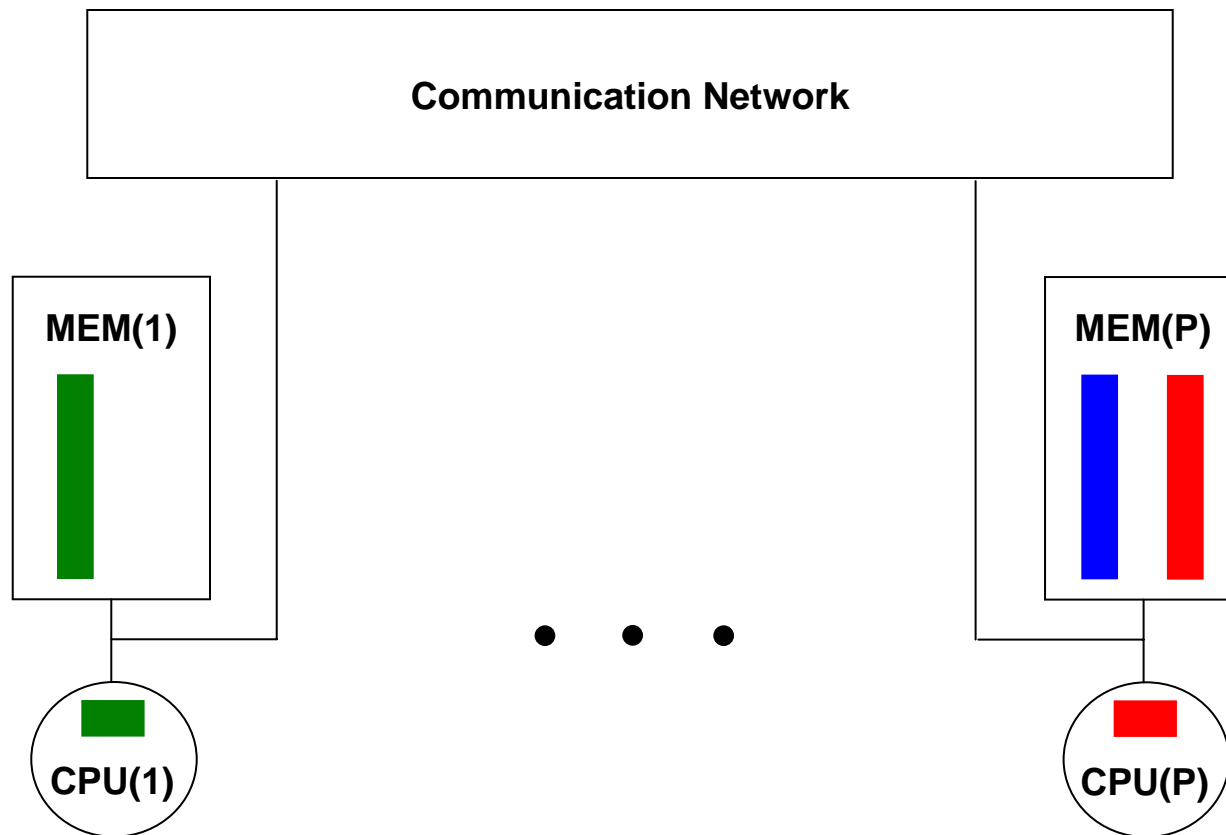
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**Mathematical Sciences Department**  
**January 18, 2005**

\* Based mostly on joint work with R. Nelson

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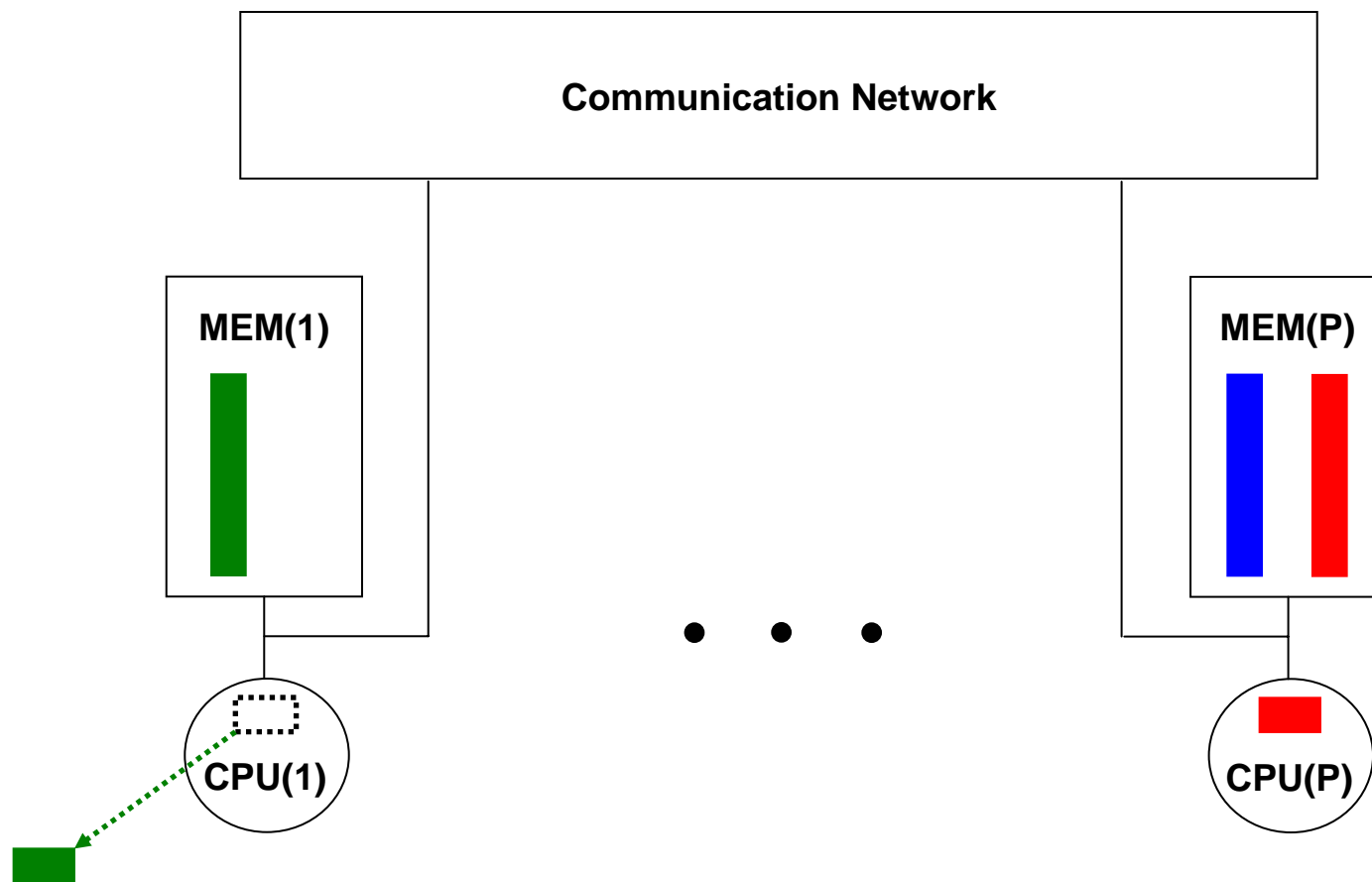
# Problem Motivation

## Parallel Computing Example



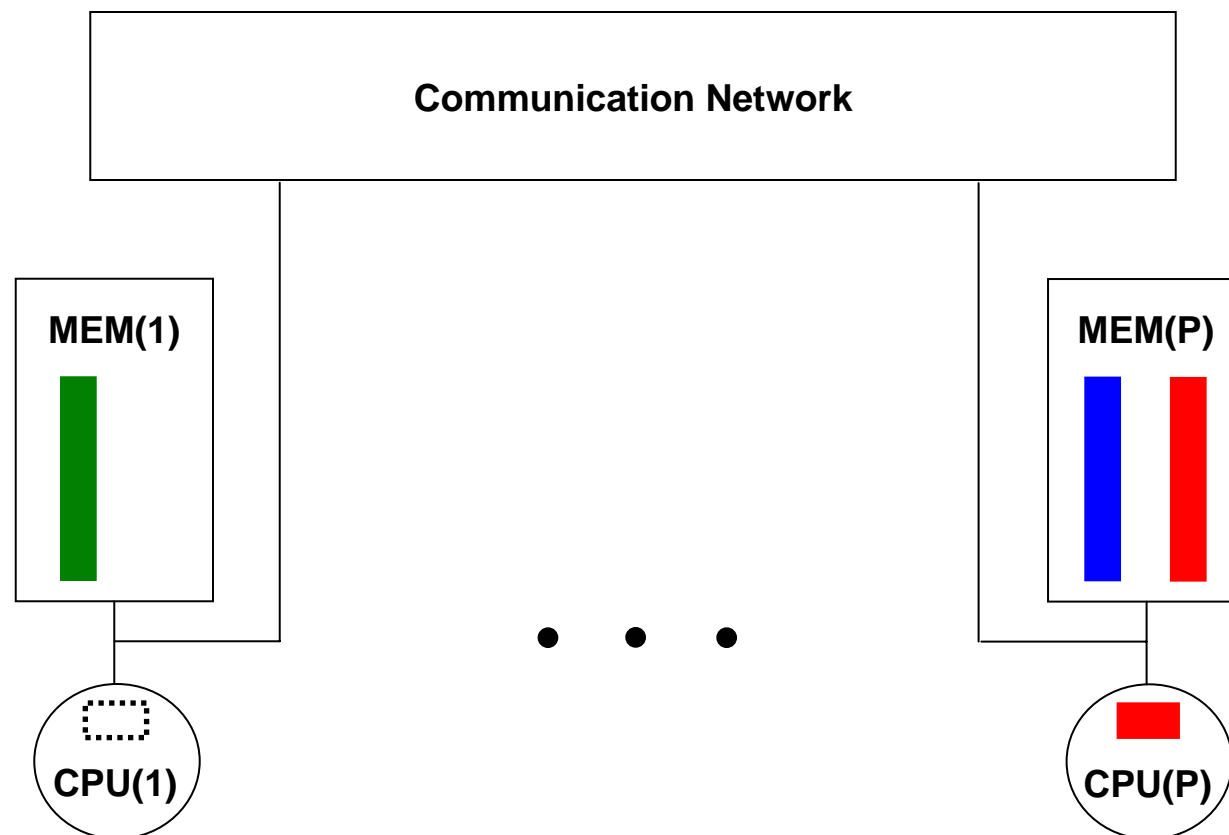
# Problem Motivation

## Parallel Computing Example



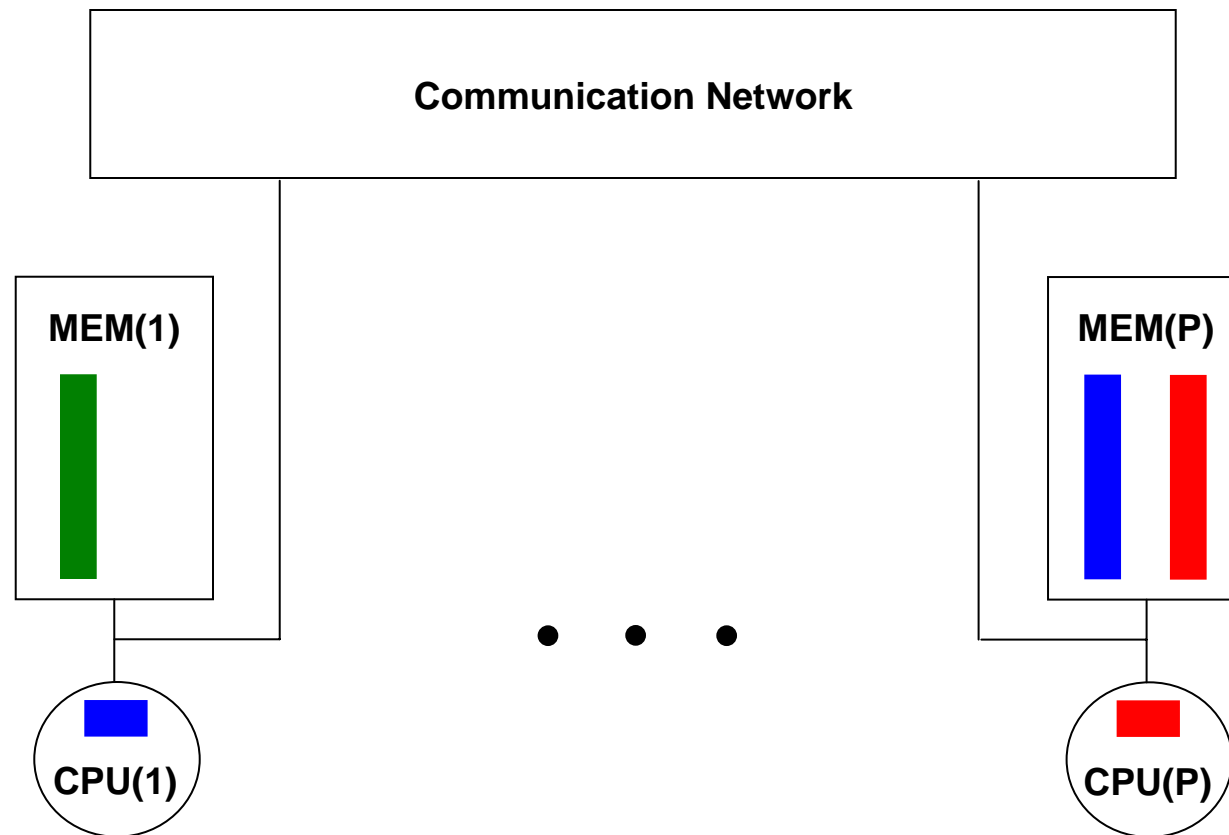
## Problem Motivation

**Dynamic Scheduling Tradeoff: Should we leave CPU(1) idle?**



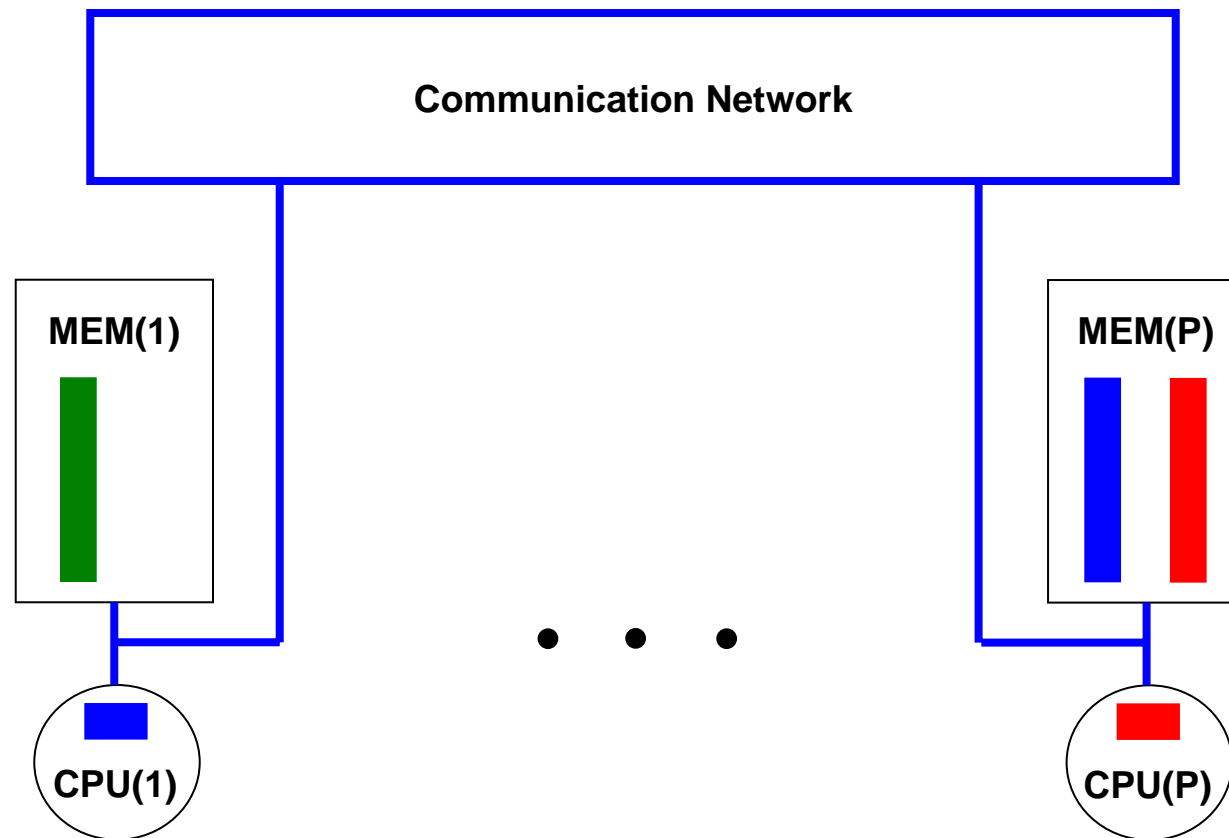
## Problem Motivation

Dynamic Scheduling Tradeoff: Or serve **Blue-Task** at CPU(1)?



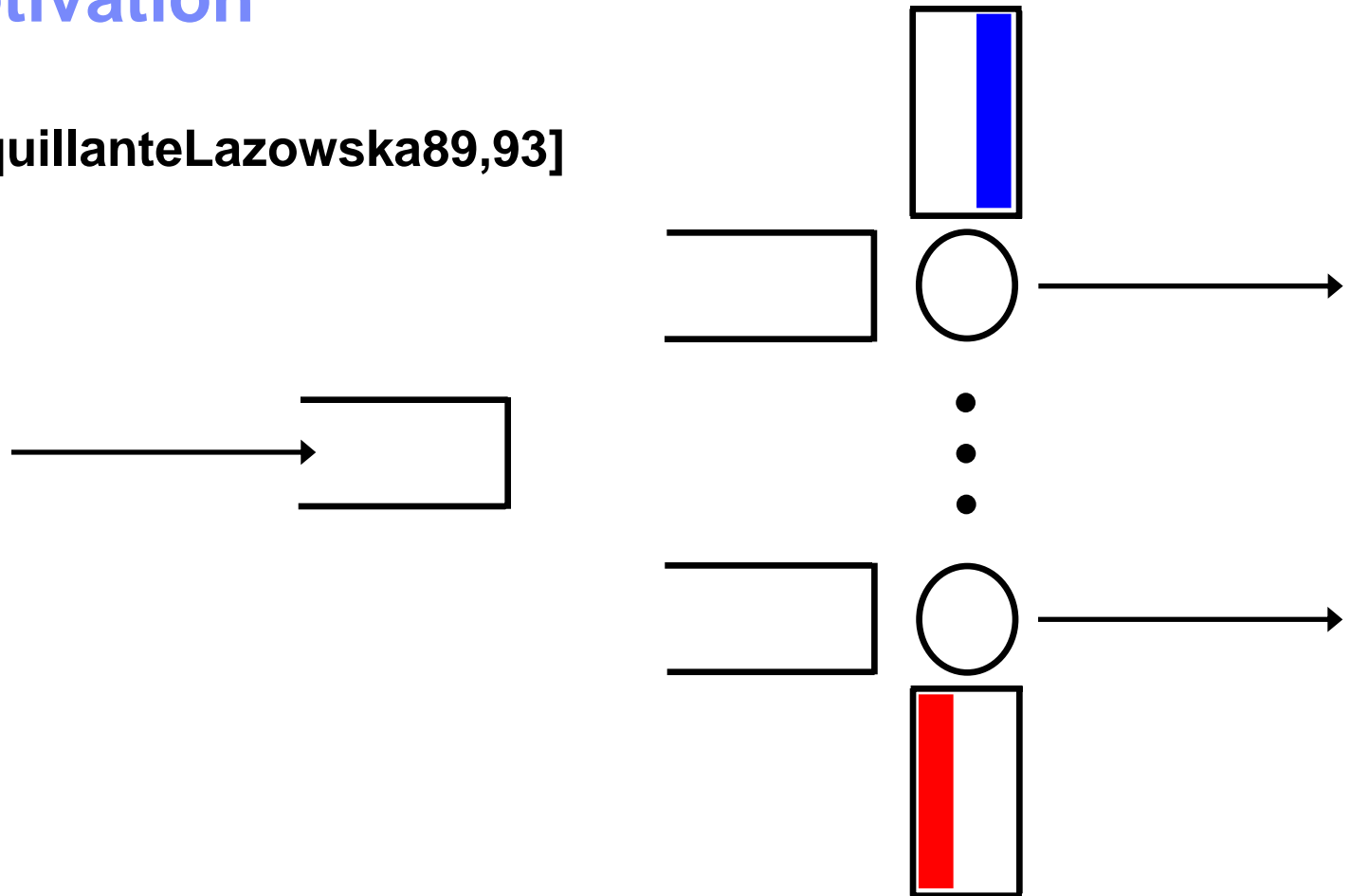
## Problem Motivation

Dynamic Scheduling Tradeoff: But larger **Blue-Task** service time



# Problem Motivation

## Cache Affinity [SquillanteLazowska89,93]



### Key Points of Fundamental Tradeoff

- Customers can be served on any server of a parallel-server queueing system
- Each customer is served most efficiently on one the servers
- Load imbalance among queues occurs due to stochastic properties of system

## General Overview

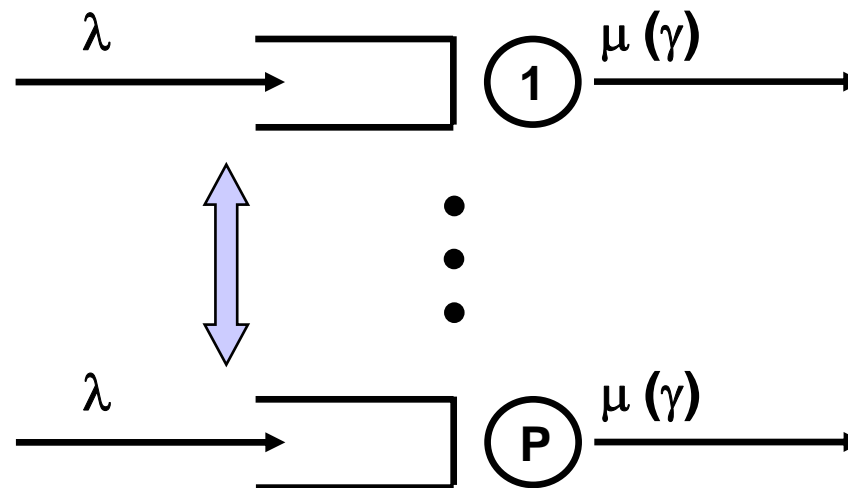
- Optimal Scheduling Policy
  - Fluid control problem:  $c\mu$ -type scheduling policy
  - Brownian control problem: dynamic threshold-type scheduling policy
  - Optimal threshold settings for dynamic scheduling policy depend upon stochastic properties and traffic intensity of system
  
- Analysis of Dynamic Threshold Scheduling
  - Consider generalized threshold scheduling policy
  - Matrix-analytic analysis and fix-point solution, asymptotically exact
  - Numerical experiments
  - Optimal settings of dynamic scheduling policy thresholds
  
- Stochastic Derivative-Free Optimization



## Scheduling Policy

$$\beta \equiv (T_r^l, T_s^l, T_r^u, T_s^u), \quad 0 \leq \{T_r^l, T_s^l\} \leq \{T_r^u, T_s^u\} < \infty$$

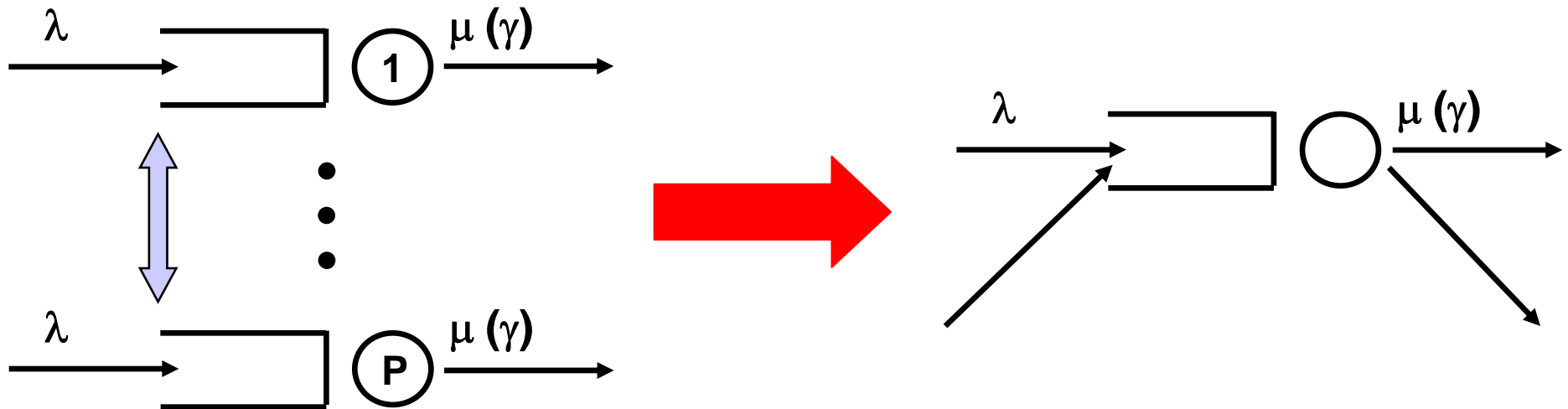
$$K \equiv \max\{T_r^l, T_s^l\}, \quad \hat{T} \equiv \max\{T_r^u, T_s^u\}$$



## Scheduling Model

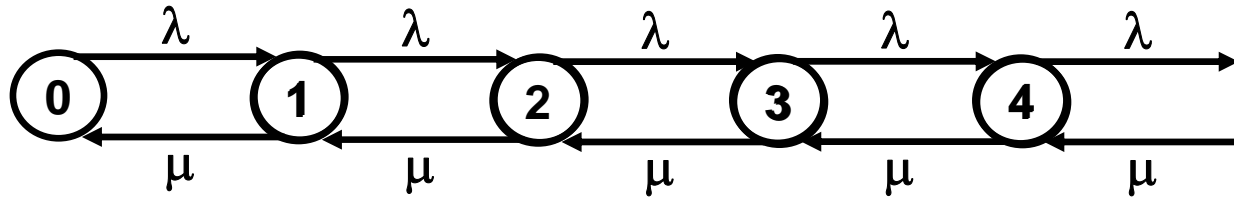
$$\beta \equiv (T_r^l, T_s^l, T_r^u, T_s^u), \quad 0 \leq \{T_r^l, T_s^l\} \leq \{T_r^u, T_s^u\} < \infty$$

$$K \equiv \max\{T_r^l, T_s^l\}, \quad \hat{T} \equiv \max\{T_r^u, T_s^u\}$$



## Background on Matrix-Analytic Methods

BD process

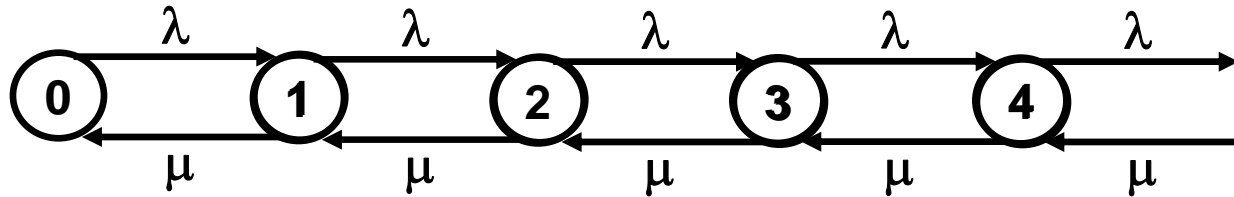


Define  $\pi_n \equiv \lim_{t \rightarrow \infty} P[N(t) = n]$ ,  $\boldsymbol{\pi} \equiv (\pi_0, \pi_1, \pi_2, \dots)$

$$Q = \begin{bmatrix} -\lambda & \lambda & 0 & 0 & 0 & \dots \\ \mu & -(\lambda + \mu) & \lambda & 0 & 0 & \dots \\ 0 & \mu & -(\lambda + \mu) & \lambda & 0 & \dots \\ 0 & 0 & \mu & -(\lambda + \mu) & \lambda & \dots \\ \vdots & \vdots & \vdots & \vdots & \vdots & \ddots \end{bmatrix}$$

## Background on Matrix-Analytic Methods

BD process



Then solution of  $\pi \mathbf{Q} = 0$  and  $\pi \mathbf{e} = 1$  given by

$$\pi_i = \pi_0 \rho^i$$

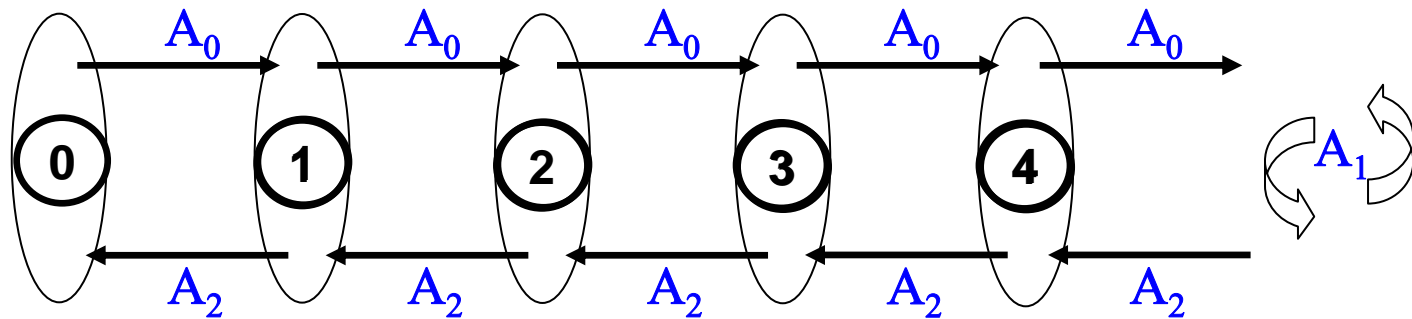
where

$$0 = \lambda - \rho(\lambda + \mu) + \rho^2 \mu, \quad \rho < 1 \quad \Rightarrow \quad \rho = \lambda / \mu < 1$$

$$\pi \mathbf{e} = 1 \quad \Rightarrow \quad \pi_0 \sum_{i=0}^{\infty} \rho^i = 1 \quad \Rightarrow \quad \pi_0 = 1 - \rho$$

# Background on Matrix-Analytic Methods

QBD process

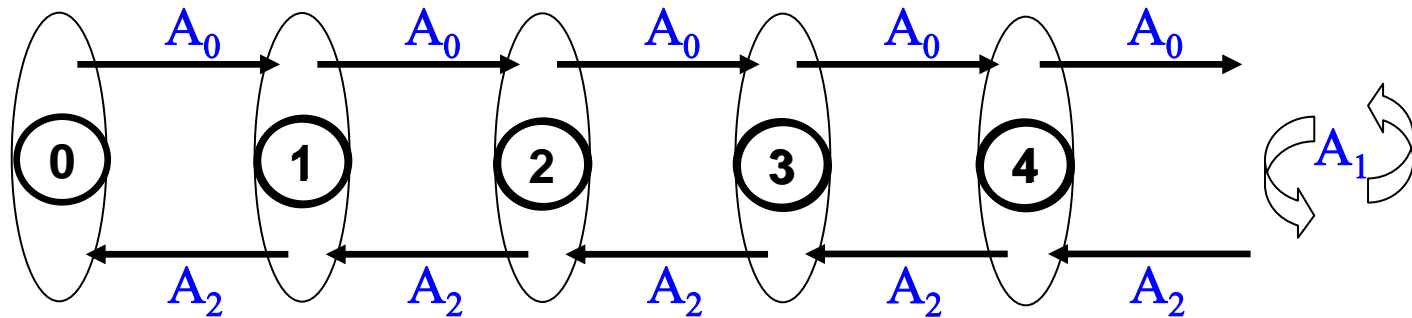


Define  $\pi_n \equiv (\pi_{n,1}, \dots, \pi_{n,m})$ ,  $\pi \equiv (\pi_0, \pi_1, \pi_2, \dots)$

$$Q = \begin{bmatrix} B_0 & A_0 & 0 & 0 & 0 & \dots \\ A_2 & A_1 & A_0 & 0 & 0 & \dots \\ 0 & A_2 & A_1 & A_0 & 0 & \dots \\ 0 & 0 & A_2 & A_1 & A_0 & \dots \\ \vdots & \vdots & \vdots & \vdots & \vdots & \ddots \end{bmatrix}$$

## Background on Matrix-Analytic Methods

QBD process



Then solution of  $\pi Q = 0$  and  $\pi e = 1$  given by

$$\pi_i = \pi_0 R^i$$

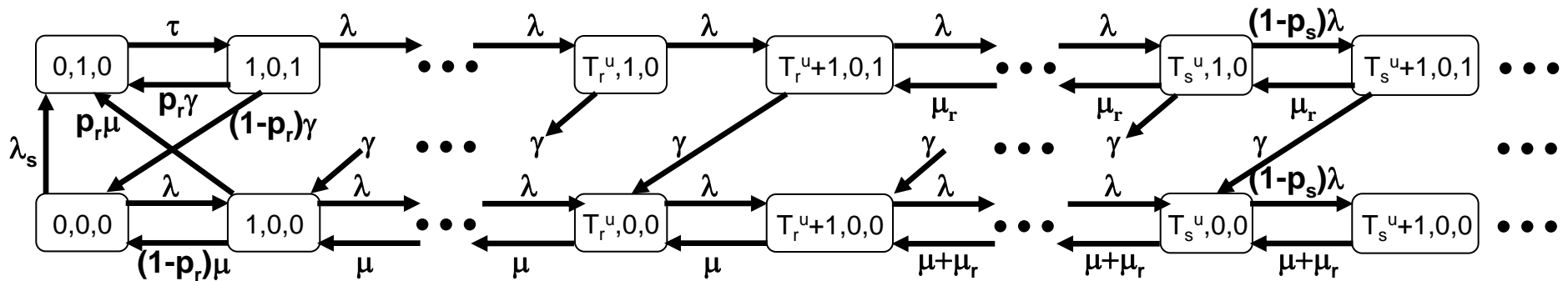
where

$$0 = A_0 + RA_1 + R^2A_2, \quad \text{sp}(R) < 1$$

$$0 = \pi_0(B_0 + RA_2), \quad \pi_0(I - R)^{-1}e = 1$$

# Mathematical Analysis

$$T_r^l = T_s^l = K = 1, \quad T_r^u < T_s^u = \hat{T}, \quad \exp(\lambda), \exp(\mu), \exp(\gamma)$$



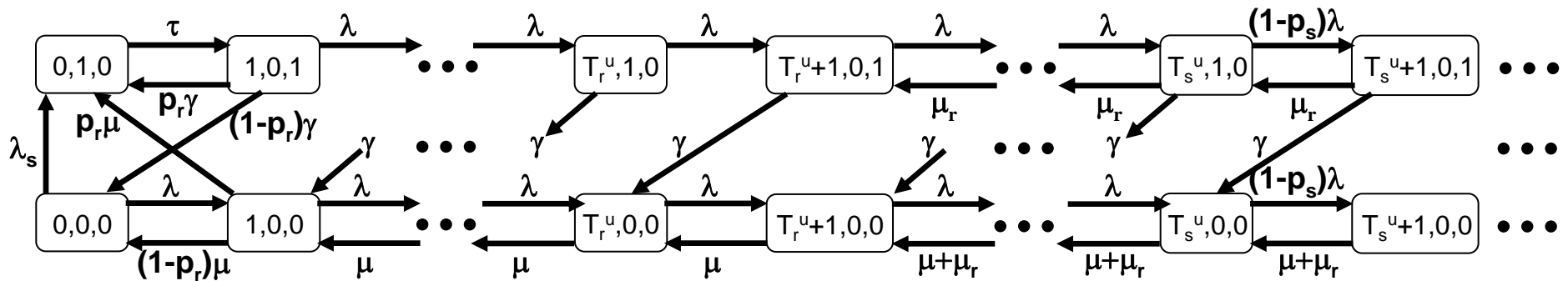
## State Vector ( i, j, v ):

- i: total number of customers waiting or receiving service at the server of interest
- j: number of customers in the process of being migrated to the server of interest
- v: K-bit binary vector denoting customer type of up to the first K customers at the server



# Mathematical Analysis

$$T_r^l = T_s^l = K = 1, \quad T_r^u < T_s^u = \hat{T}, \quad \exp(\lambda), \exp(\mu), \exp(\gamma)$$



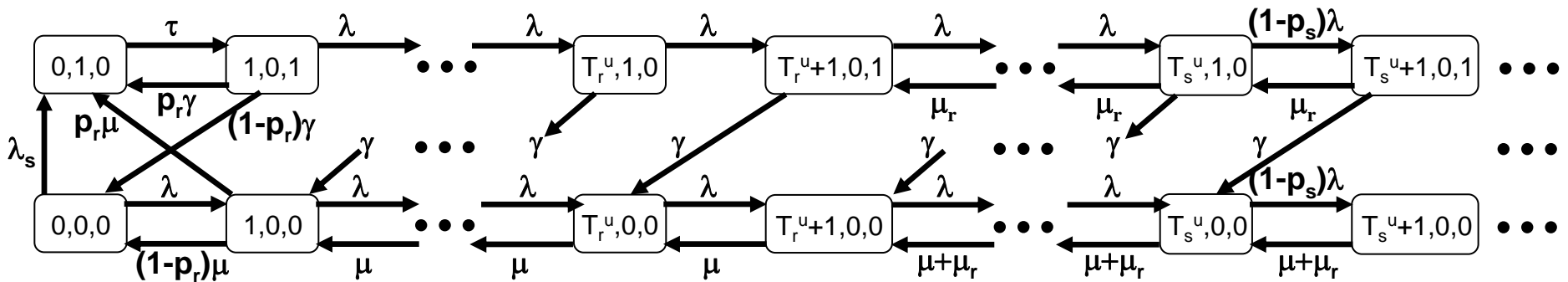
Define  $\pi_0 \equiv (\pi_{0,0,0}, \pi_{0,1,0})$ ,  $\pi_n \equiv (\pi_{n,0,0}, \pi_{n,0,1})$ ,  $n > 0$

$$Q = \begin{bmatrix} B_{00} & B_{01} & 0 & 0 & 0 & \dots \\ B_{10} & B_{11} & A_0 & 0 & 0 & \dots \\ 0 & A_2 & A_1 & A_0 & 0 & \dots \\ 0 & 0 & A_2 & A_1 & A_0 & \dots \\ \vdots & \vdots & \vdots & \vdots & \vdots & \ddots \end{bmatrix}$$



# Mathematical Analysis

$$T_r^l = T_s^l = K = 1, \quad T_r^u < T_s^u = \hat{T}, \quad \exp(\lambda), \exp(\mu), \exp(\gamma)$$



Then solution of  $\pi Q = 0, \pi e = 1$  is given by

$$\pi_{\hat{T}+k} = \pi_{\hat{T}} \mathbf{R}^k, \quad k \geq 0$$

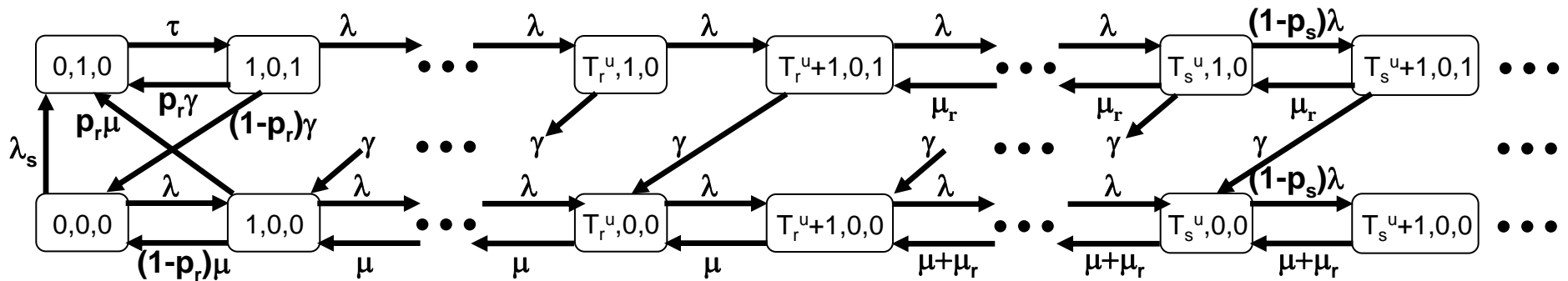
$$0 = \mathbf{A}_0 + \mathbf{R}\mathbf{A}_1 + \mathbf{R}^2\mathbf{A}_2, \quad \text{sp}(\mathbf{R}) < 1$$

$$0 = (\pi_0, \pi_1, \dots, \pi_{\hat{T}}) \begin{bmatrix} \mathbf{B}_{00} & \mathbf{B}_{01} \\ \mathbf{B}_{10} & \mathbf{B}_{11} + \mathbf{R}\mathbf{A}_2 \end{bmatrix}$$

$$1 = (\pi_0, \pi_1, \dots, \pi_{\hat{T}-1})\mathbf{e} + \pi_{\hat{T}}(\mathbf{I} - \mathbf{R})^{-1}\mathbf{e}$$

# Mathematical Analysis

$$T_r^l = T_s^l = K = 1, \quad T_r^u < T_s^u = \hat{T}, \quad \exp(\lambda), \exp(\mu), \exp(\gamma)$$



Probabilistic Interpretation of  $\mathbf{R}$ :

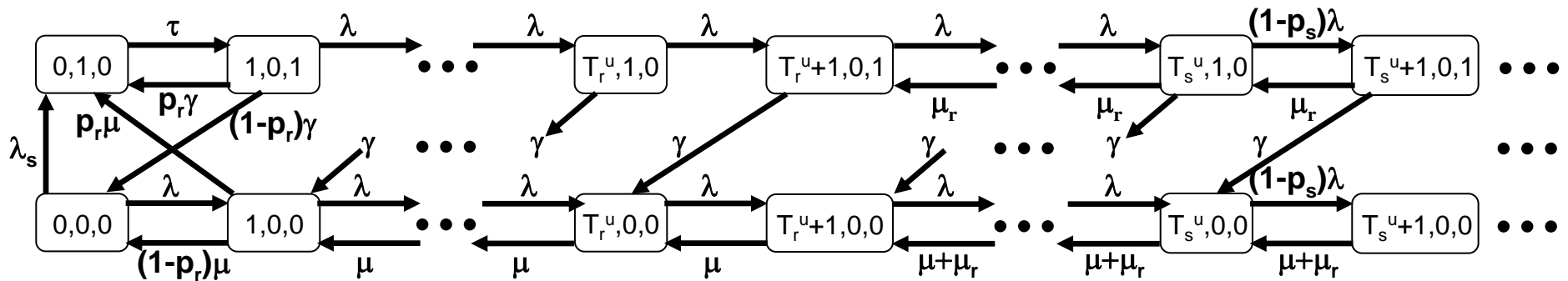
$$r_{00} = \frac{\lambda(1 - p_s)}{\mu + \mu_r}, \quad r_{01} = 0$$

$$r_{11} = \frac{\lambda(1 - p_s) + \gamma + \mu_r - ((\lambda(1 - p_s) + \gamma + \mu_r)^2 - 4\lambda(1 - p_s)\mu_r)^{1/2}}{2\mu_r}$$

$$r_{10} = \frac{r_{11}^2 \gamma}{(\mu + \mu_r)(1 - r_{11})}$$

# Mathematical Analysis

$$T_r^\ell = T_s^\ell = K = 1, \quad T_r^u < T_s^u = \hat{T}, \quad \exp(\lambda), \exp(\mu), \exp(\gamma)$$



Let  $I, J, \underline{V}$  be generic r.v.s of the process

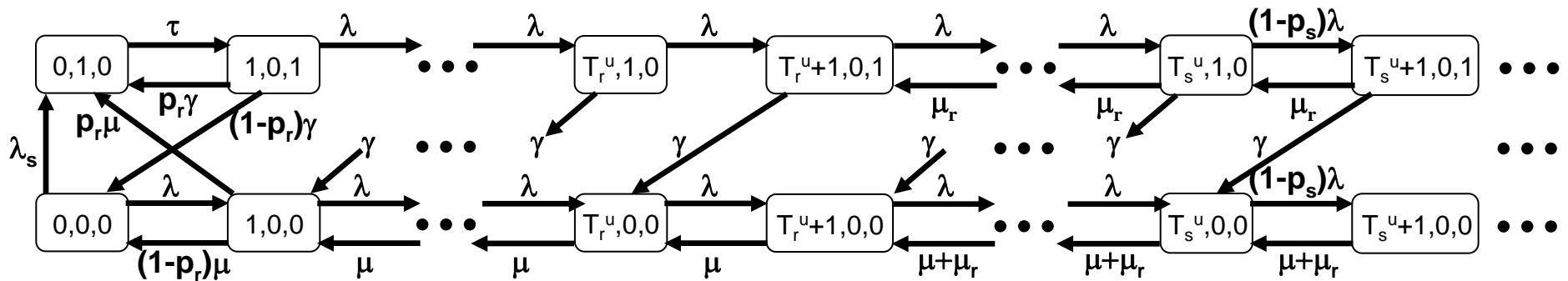
$$\begin{aligned}
 p_s &= \mathbf{P} \left[ \text{Find Queue with } I + J < T_s^\ell \text{ in } \leq L_p \text{ Random Probes} \right] \\
 &= 1 - \left( 1 - \mathbf{P} \left[ I + J < T_s^\ell \right] \right)^{L_p}
 \end{aligned}$$

Equating inflow and outflow of SI migrated customers:

$$\lambda_s \mathbf{P} \left[ I + J < T_s^\ell \right] = \lambda p_s \mathbf{P} \left[ I \geq T_s^u \right]$$

# Mathematical Analysis

$$T_r^l = T_s^l = K = 1, \quad T_r^u < T_s^u = \hat{T}, \quad \exp(\lambda), \exp(\mu), \exp(\gamma)$$



Let  $I, J, \underline{V}$  be generic r.v.s of the process

$$p_r = \text{P} [\text{Find Queue with } I > T_r^u \text{ in } \leq L_p \text{ Random Probes}]$$

$$= 1 - (1 - \text{P} [I > T_r^u])^{L_p}$$

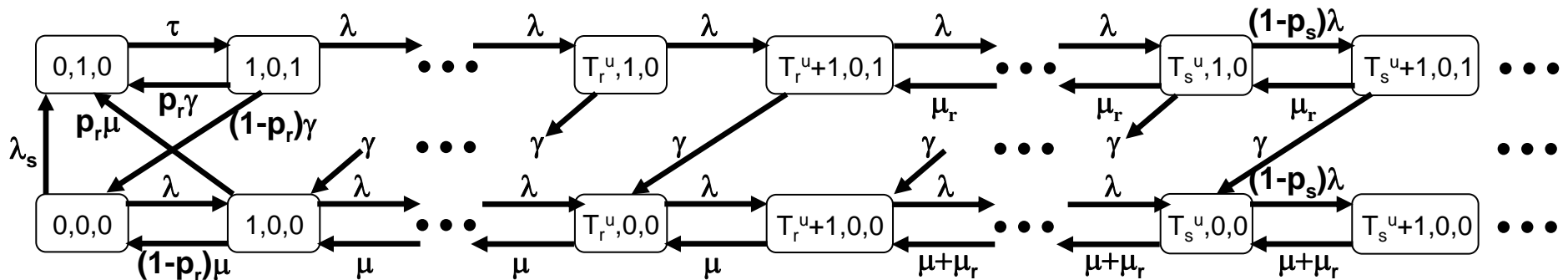
Equating inflow and outflow of RI migrated customers:

$$\mu_r \text{P} [I > T_r^u] = \mu p_r \text{P} [I + J \leq T_r^l, I > 0, v_1 = 0] +$$

$$\gamma p_r \text{P} [I + J \leq T_r^l, I > 0, v_1 = 1]$$

# Mathematical Analysis

$$T_r^l = T_s^l = K = 1, \quad T_r^u < T_s^u = \hat{T}, \quad \exp(\lambda), \exp(\mu), \exp(\gamma)$$



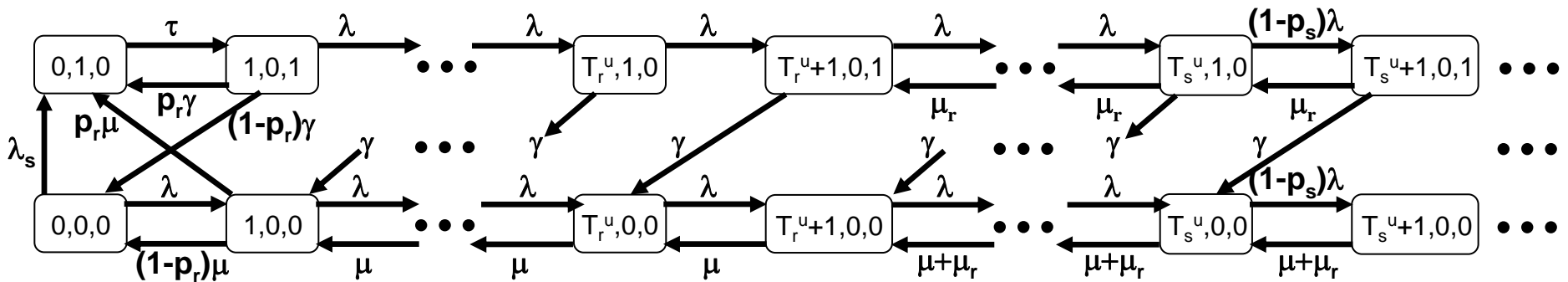
## Final Solution Obtained via Fix-Point Iteration

1. Initialize  $(\mathbf{p}_s, \lambda_s, \mathbf{p}_r, \mu_r)$
2. Compute stationary probability vector in terms of  $(\mathbf{p}_s, \lambda_s, \mathbf{p}_r, \mu_r)$
3. Compute new values of  $(\mathbf{p}_s, \lambda_s, \mathbf{p}_r, \mu_r)$  in terms of stationary vector
4. Goto 2 until differences between iteration values are arbitrarily small



# Mathematical Analysis

$$T_r^l = T_s^l = K = 1, \quad T_r^u < T_s^u = \hat{T}, \quad \exp(\lambda), \exp(\mu), \exp(\gamma)$$



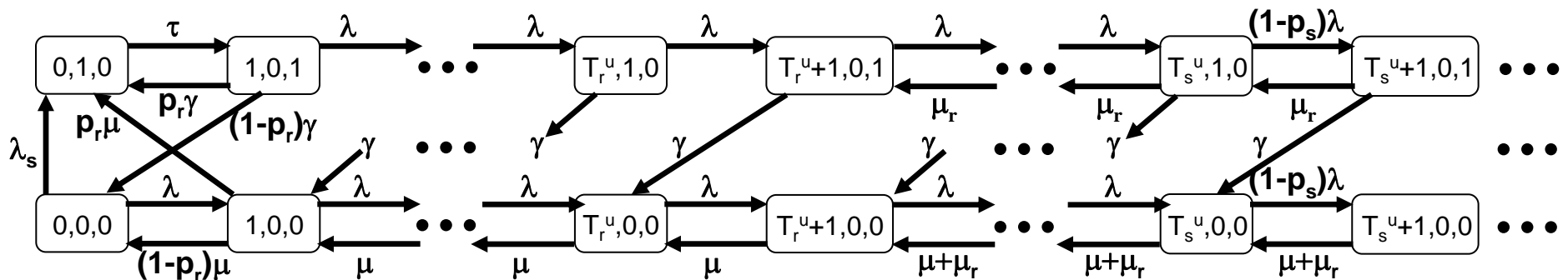
Stability criterion:  $sp(\mathbf{R}) = r_{00} < 1$ ,  $\mathbf{A} \equiv \mathbf{A}_0 + \mathbf{A}_1 + \mathbf{A}_2$  reducible

$$E[\mathcal{T}] = \lambda^{-1} \left\{ \sum_{k=1}^{\hat{T}-1} k \pi_k \mathbf{e} + \hat{T} \pi_{\hat{T}} (\mathbf{I} - \mathbf{R})^{-1} \mathbf{e} + \pi_{\hat{T}} (\mathbf{I} - \mathbf{R})^{-2} \mathbf{R} \mathbf{e} \right\}$$

$$p_{mt} = (\pi_1, \dots, \pi_K) \nu + \sum_{k=K+1}^{\hat{T}-1} \pi_k \nu' + \pi_{\hat{T}} (\mathbf{I} - \mathbf{R})^{-1} \nu'$$

# Mathematical Analysis

$$T_r^l = T_s^l = K = 1, \quad T_r^u < T_s^u = \hat{T}, \quad \exp(\lambda), \exp(\mu), \exp(\gamma)$$



Tail asymptotics:

$$P[Q > x] = \frac{\pi_{\hat{T}}^{\mathbf{v}}}{1 - \eta} \eta^x + o(\eta^x), \quad \text{as } x \rightarrow \infty,$$

$$P[Q > x] \sim \frac{\pi_{\hat{T}}^{\mathbf{v}}}{1 - \eta} \eta^x, \quad \text{as } x \rightarrow \infty,$$

where  $\mathbf{u}$  and  $\mathbf{v}$  are the left and right eigenvectors corresponding to  $\eta = r_{00}$  normalized by  $\mathbf{u}\mathbf{e} = 1$  and  $\mathbf{u}\mathbf{v} = 1$

## Mathematical Analysis

### General K, Pure Sender-Initiated Policy

Probabilistic Interpretation of  $\mathbf{R}$ :

$$\begin{aligned}
 r_{00} &= \lambda(1 - p_s)/\mu \\
 r_{nm} &= 0, \quad 0 \leq n < m \leq 2^K - 1 \quad (\text{lower triangular}) \\
 r_{nn}^0 &= \Phi(n), \quad 1 \leq n \leq 2^K - 1 \\
 r_{nm}^0 &= C_k \Phi(n)^{k+1} \Psi(n)^k, \quad (*) \\
 &\quad n > m > 0 \text{ s.t. } b_m = \mathcal{D}_k(b_n), \quad 1 \leq k < K \\
 r_{n0} &= \frac{\gamma \sum_{k=1}^n r_{nk} r_{k1} + \mu \sum_{k=1}^{n-1} r_{nk} r_{k0}}{\lambda(1 - p_s) + \mu(1 - r_{00} - r_{nn})}, \\
 &\quad 1 \leq n \leq 2^K - 1
 \end{aligned}$$



# Mathematical Analysis

## General K

Departure Transitions from State $(i, j, \underline{v})$ , $i > 0$ , $j \leq (K - i)^+$	
To State	Rate
$(i - 1, j, \mathcal{D}_1(\underline{v})), i + j \leq T_r^\ell \leq K$	$(1 - p_r)e(\underline{v})$
$(i - 1, j + 1, \mathcal{D}_1(\underline{v})), i + j \leq T_r^\ell \leq K$	$p_r e(\underline{v})$
$(i - 1, j, \mathcal{D}_1(\underline{v})), T_r^\ell < i + j \leq T_r^u$	$e(\underline{v})$
$(i - 1, j, \underline{v}), \underline{v} \neq 0, i > T_r^u$	$\mu_r$
$(i - 1, j, \mathcal{D}_1(\underline{v})), \underline{v} \neq 0, i > T_r^u$	$e(\underline{v})$
$(i - 1, j, \mathcal{D}_1(\underline{v})), \underline{v} = 0, i > T_r^u$	$e(\underline{v}) + \mu_r$

$$\mathcal{D}_k(\underline{v}) = (v_{k+1}, \dots, v_K, \underbrace{0, \dots, 0}_{k \text{ times}})$$

# Mathematical Analysis

## General K

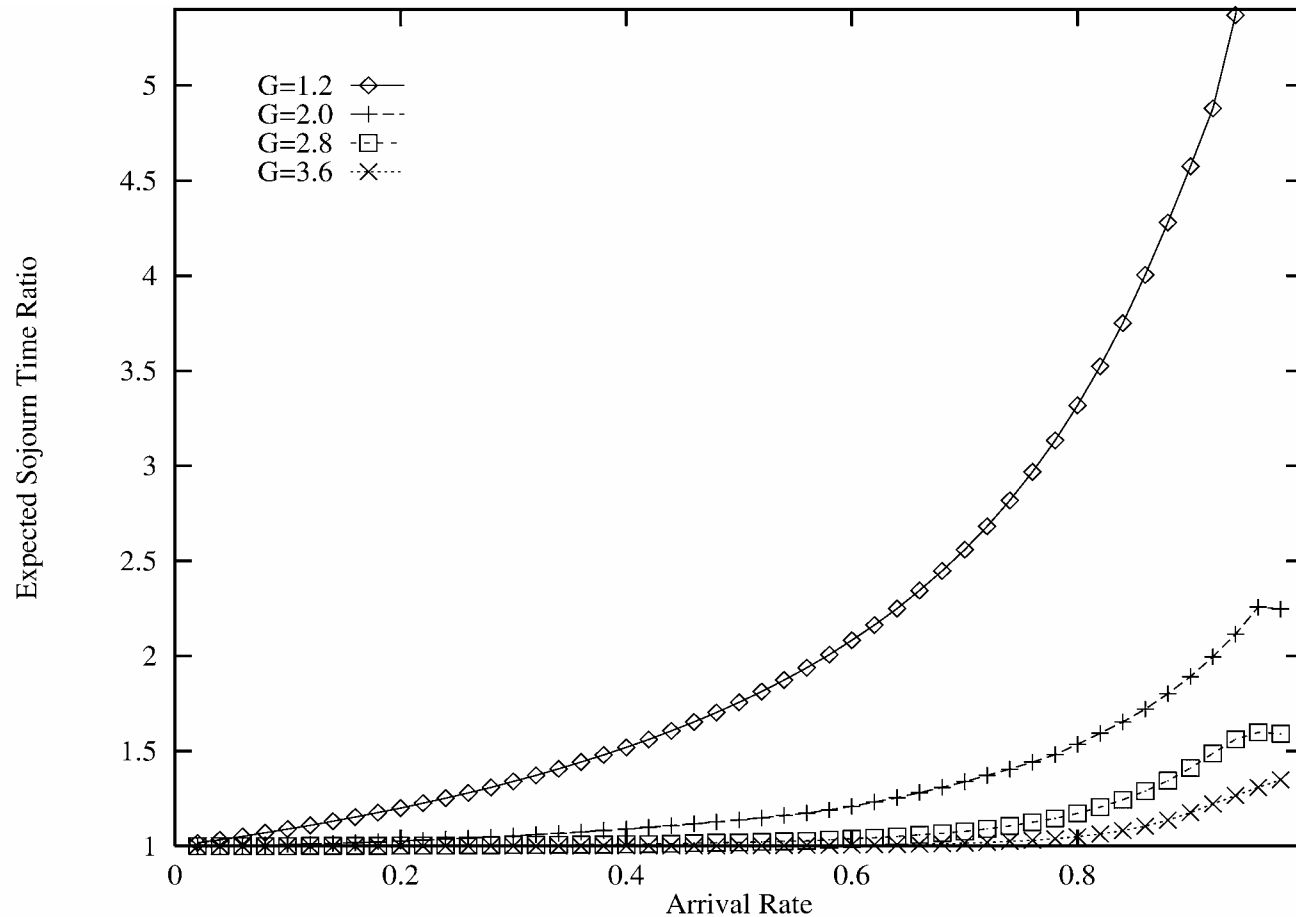
Arrival Transitions from State $(i, j, \underline{v})$ , $i \geq 0$ , $j \leq (K - i)^+$	
To State	Rate
$(i + 1, j - 1, \mathcal{A}_{i+1}(\underline{v}))$ , $j > 0$	$j\theta$
$(i, j + 1, \underline{v})$ , $i + j < T_s^l \leq K$	$\lambda_s$
$(i + 1, j, \underline{v})$ , $i < K$ , $i + j < K$	$\lambda$
$(i + 1, j, \underline{v})$ , $K \leq i < T_s^u$	$\lambda$
$(i + 1, j, \underline{v})$ , $i \geq T_s^u$	$\lambda(1 - p_s)$

$$\mathcal{A}_i(\underline{v}) = (v_1, v_2, \dots, v_{i-1}, 1, \underbrace{0, 0, \dots, 0}_{K-i \text{ times}}).$$

# Numerical Results

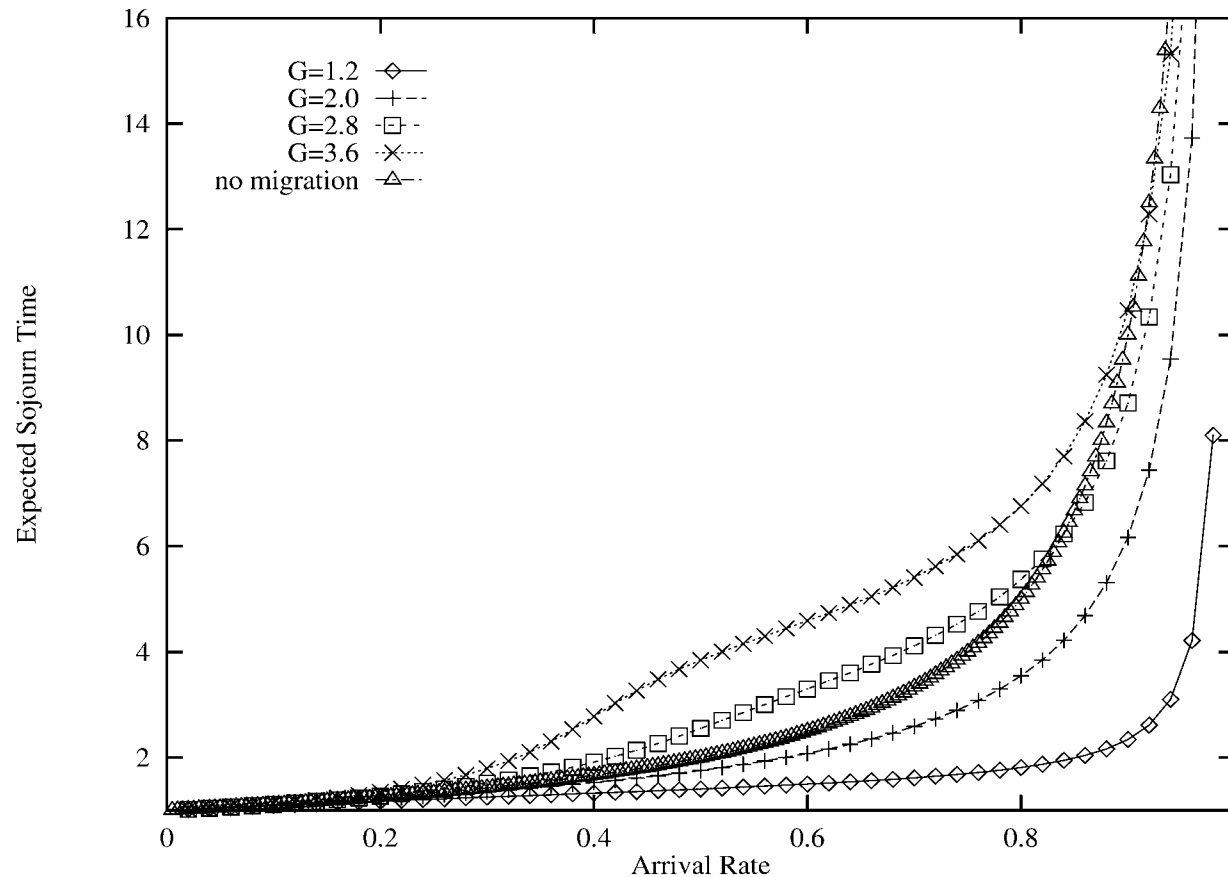
Optimal Threshold Values												
$\lambda$	$1/\gamma = 1.6$				$1/\gamma = 2.4$				$1/\gamma = 3.2$			
	$\hat{T}_r^l$	$\hat{T}_s^l$	$\hat{T}_r^u$	$\hat{T}_s^u$	$\hat{T}_r^l$	$\hat{T}_s^l$	$\hat{T}_r^u$	$\hat{T}_s^u$	$\hat{T}_r^l$	$\hat{T}_s^l$	$\hat{T}_r^u$	$\hat{T}_s^u$
0.04	1	1	1	1	0	1	-	2	0	1	-	3
0.08	1	1	1	1	1	1	2	2	1	1	3	3
0.12	1	1	1	1	1	1	2	2	1	1	4	4
0.16	1	1	1	1	1	1	2	2	1	1	4	4
0.20	1	1	1	1	1	1	3	3	1	1	4	4
0.24	1	1	1	1	1	1	3	3	1	1	5	5
0.28	1	1	1	1	1	1	3	3	1	1	5	5
0.32	1	1	1	1	1	1	3	3	1	1	5	5
0.36	1	1	1	1	1	1	3	3	1	1	5	5
0.40	1	1	1	1	1	1	3	3	1	1	6	6
0.44	1	1	1	1	1	1	3	3	1	1	6	6
0.48	1	1	1	1	1	1	4	4	1	1	6	6
0.52	1	1	1	1	1	1	4	4	1	1	7	7
0.56	1	1	1	1	1	1	4	4	1	1	7	7
0.60	1	1	1	1	1	1	4	4	1	1	8	8
0.64	1	1	1	1	1	1	4	4	1	1	8	8
0.68	1	1	2	2	1	1	4	4	1	1	8	8
0.72	1	1	2	2	1	1	5	5	1	1	9	9
0.76	1	1	2	2	1	1	5	5	1	1	9	9
0.80	1	1	2	2	1	1	5	5	1	1	10	10
0.84	1	1	3	3	1	1	6	6	1	1	10	10
0.88	1	1	3	3	1	1	7	7	1	1	12	12
0.92	1	1	4	4	1	1	9	9	1	1	12	12
0.96	1	1	12	12	1	1	12	12	1	1	12	12
0.98	1	1	12	12	1	1	12	12	1	1	12	12

## Numerical Results



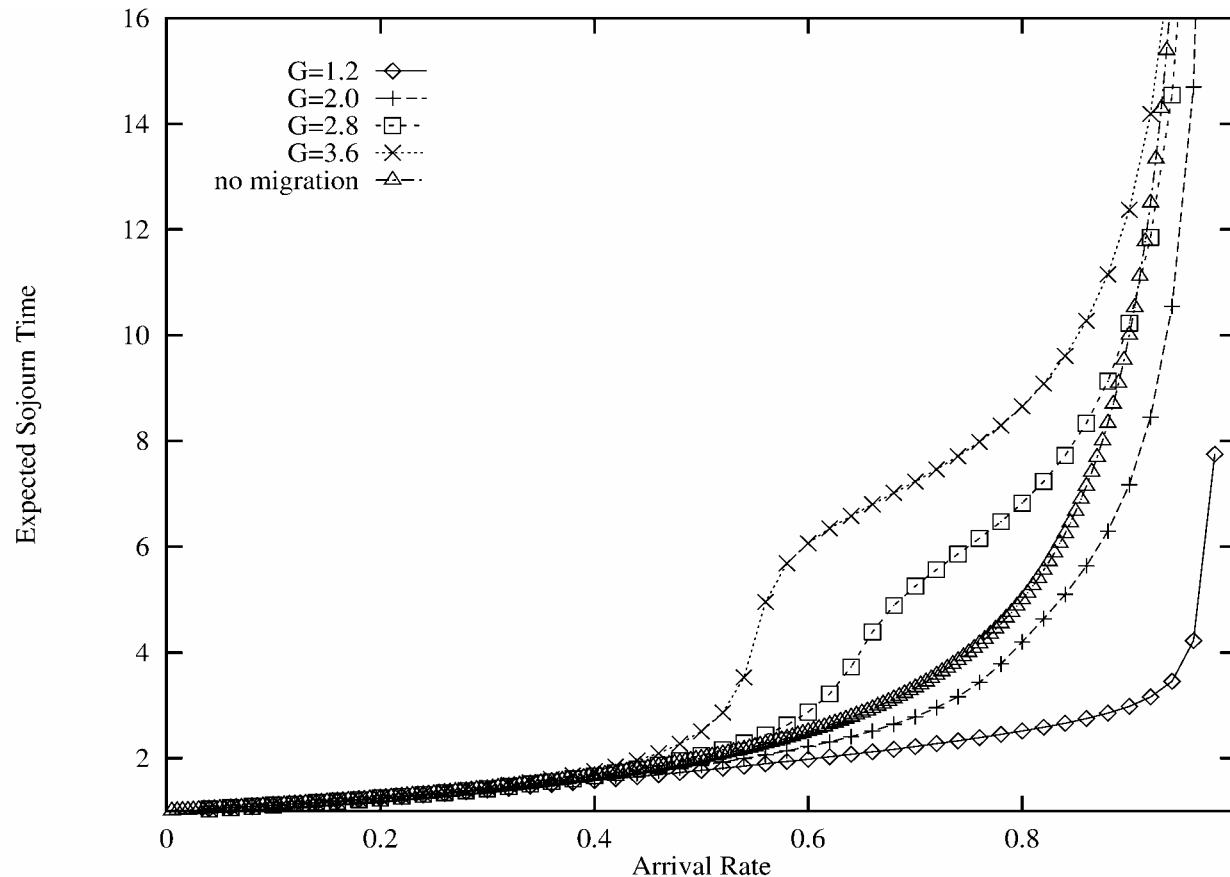
Relative Performance Benefits of the Migratory Scheduling Policy with Optimal Thresholds over the No-Migration Policy

# Numerical Results



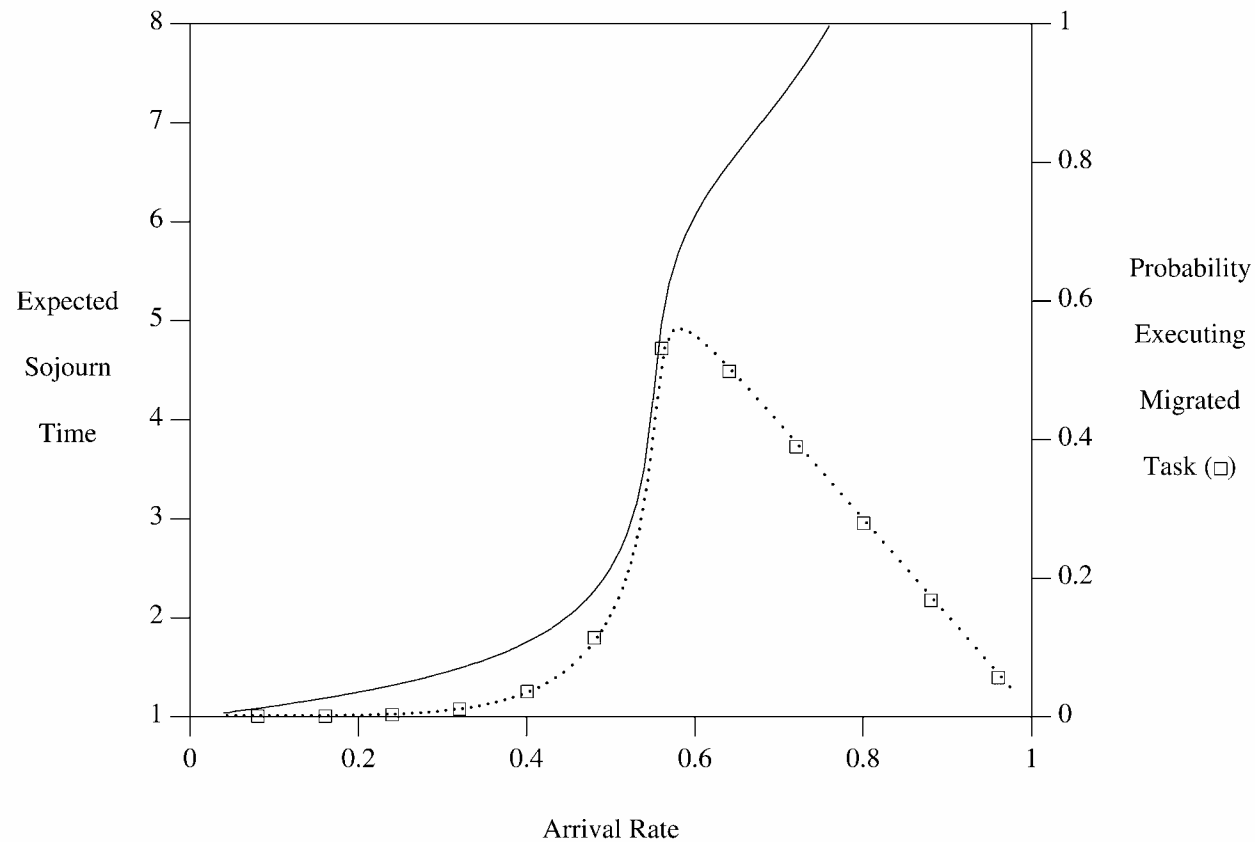
Expected Sojourn Times for the No-Migration Policy and the Migratory Scheduling Policy with Thresholds  $\beta = (T_r^l, T_s^l, T_r^u, T_s^u) = (1, 1, 2, 2)$

# Numerical Results



Expected Sojourn Times for the No-Migration Policy and the Migratory Scheduling Policy with Thresholds  $\beta = (T_r^l, T_s^l, T_r^u, T_s^u) = (2, 2, 4, 4)$

# Numerical Results



Expected Sojourn Time, and Probability of Executing a Migrated Task, for the Migratory Scheduling Policy with Thresholds  $\beta = (T_r^l, T_s^l, T_r^u, T_s^u) = (2, 2, 4, 4)$

# Stochastic Derivative-Free Optimization

- Internal Model
- External Model
- Trust Region



## General Overview

- Optimal Scheduling Policy
  - Fluid control problem:  $c\mu$  type scheduling policy
  - Brownian control problem: dynamic threshold-type scheduling policy
  
- Analysis of Dynamic Threshold Scheduling
  - Consider generalized threshold scheduling policy
  - Matrix-analytic analysis and fix-point solution, asymptotically exact
  - Numerical experiments
  - Optimal settings of dynamic scheduling policy thresholds
  
- Stochastic Derivative-Free Optimization